Grid Classes and Partial Well-Order

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Outline

Pattern Containment

- A permutation $\tau = \tau(1)\cdots\tau(k)$ is contained in the permutation $\sigma = \sigma(1)\sigma(2)\cdots\sigma(n)$ if there exists a subsequence $\sigma(i_1)\sigma(i_2)\cdots\sigma(i_k)$ order isomorphic to τ.

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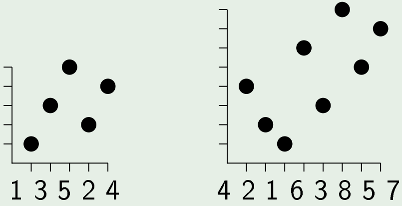

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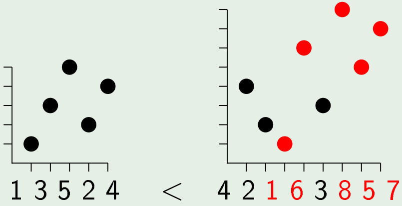

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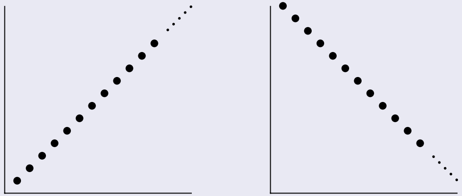
- Downsets of permutations in this partial order form permutation classes.
  i.e. $\pi \in \mathcal{C}$ and $\sigma \leq \pi$ implies $\sigma \in \mathcal{C}$.
- Typical description: basis is the set of minimal excluded elements.
  $\mathcal{C} = \mathrm{Av}(B) = \{\pi : \beta \not\leq \pi \text{ for all } \beta \in B\}$.

Cherry-picked Examples

- $\mathrm{Av}(21) = \{1, 12, 123, 1234, \ldots\}$, the increasing permutations.
- $\mathrm{Av}(12) = \{1, 21, 321, 4321, \ldots\}$, the decreasing permutations.

Typical Elements



Cherry-picked Examples

- $\oplus 21 = \text{Av}(321, 312, 231) = \{1, 12, 21, 123, 132, 213, \ldots\}$.
- $\ominus 12 = \text{Av}(123, 213, 132) = \{1, 12, 21, 231, 312, 321, \ldots\}$.

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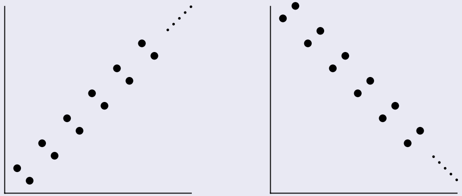

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- Graphs: containment as a graph minor. Downsets are minor-closed classes. e.g. planar graphs.
- Tournaments, words, posets all have equivalent concepts.

Antichains

- Set of pairwise incomparable permutations.

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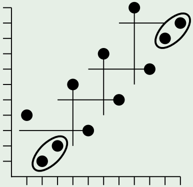
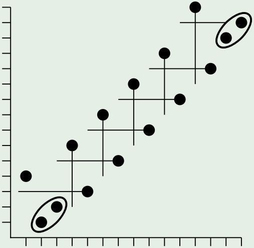
Example (Increasing Oscillating Antichain)





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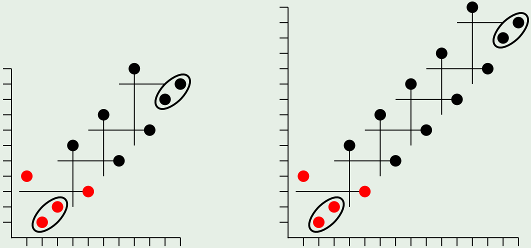

- Bottom copies of 4123 must match up: the anchor.

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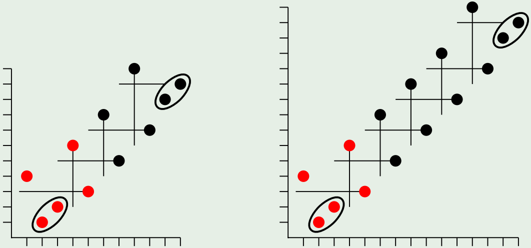

- Each point is matched in turn.

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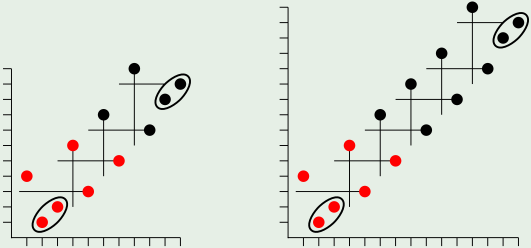

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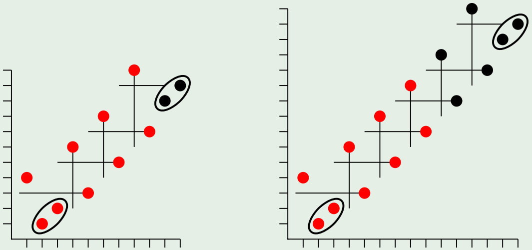

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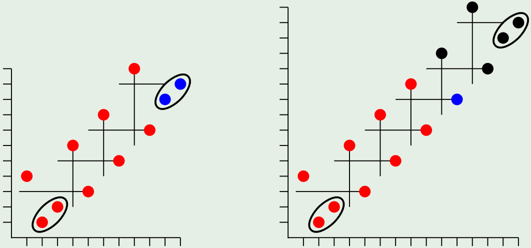

- Last pair cannot be embedded.

When are there antichains?

No infinite antichains.

- Words over a finite alphabet [Higman].
- Graphs closed under minors [Robertson and Seymour].

Infinite antichains.

- Graphs closed under induced subgraphs (or merely subgraphs). e.g. $C_3, C_4, C_5, \ldots$
- Permutations closed under containment.
- Tournaments, digraphs, ...

Partial Well Order

- A permutation class is partially well-ordered (pwo) if it contains no infinite antichains.

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Question

Can we decide whether a permutation class given by a finite basis is pwo?

- To prove pwo — Higman's theorem is useful.
- To prove not pwo — find an antichain.

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Question

Can we decide whether a hereditary property given by a finite basis is wqo?

- To prove pwo — Higman's theorem is useful.
- To prove not pwo — find an antichain.
- Other structures: well quasi-order, not pwo, but same idea.

Off topic: subgraph ordering

Theorem (Ding)

A class of graphs closed under taking subgraphs is well quasi-ordered if and only if it contains all but finitely many elements of $\{C_3, C_4, \ldots\}$ *and* $\{H_1, H_2, \ldots\}$.

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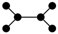
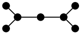
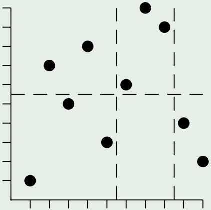
- $H_1 =$ , $H_2 =$ , $H_3 =$ , ...

Grid Classes

- Matrix $\mathcal{M}$ whose entries are permutation classes.
- $\mathrm{Grid}(\mathcal{M})$ the grid class of $\mathcal{M}$: all permutations which can be "gridded" so each cell satisfies constraints of $\mathcal{M}$.

Example

- Let $\mathcal{M} = \begin{pmatrix} \mathrm{Av}(21) & \mathrm{Av}(231) & \emptyset \\ \mathrm{Av}(123) & \emptyset & \mathrm{Av}(12) \end{pmatrix}$.



$\in \mathrm{Grid}(\mathcal{M})$

Grid classes are useful

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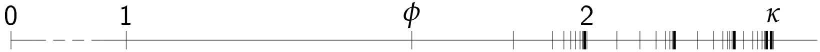
- Grid classes give a complete answer to permitted growth rates below $\kappa \approx 2.20557$ [Vatter]:



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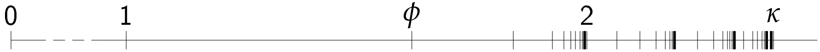


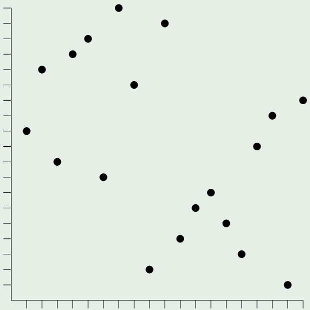
- Cf "canonical properties" of graphs [Balogh, Bollobás and Weinreich].

Monotone Grid Classes

- Special case: all cells of $\mathcal{M}$ are $\mathrm{Av}(21)$ or $\mathrm{Av}(12)$.
- Rewrite $\mathcal{M}$ as a matrix with entries in $\{0, 1, -1\}$.

Example

$$\mathcal{M} = \begin{pmatrix} 1 & 1 & 0 \\ -1 & 0 & 1 \\ 0 & 1 & -1 \end{pmatrix}$$



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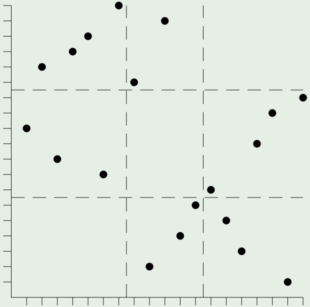

The Graph of a Matrix

- Graph of a matrix, $G(\mathcal{M})$, formed by connecting together all non-zero entries that share a row or column and are not "separated" by any other nonzero entry.

Example

$$\begin{pmatrix} \mathcal{C} & 0 & 0 & \mathcal{D} \\ 0 & 0 & \mathcal{E} & 0 \\ \mathcal{D} & \mathcal{E} & 0 & \mathcal{C} \\ 0 & 0 & 0 & \mathcal{D} \end{pmatrix}$$

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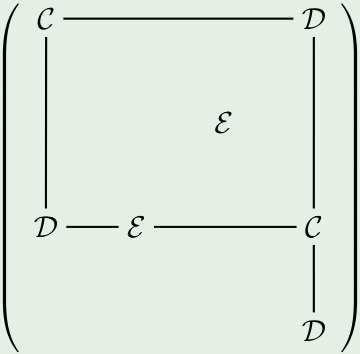
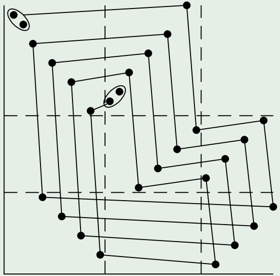


Monotone Grids and Partial Well Order

Theorem (Murphy and Vatter, 2003)

The monotone grid class $\mathrm{Grid}(\mathcal{M})$ *is pwo if and only if* $G(\mathcal{M})$ *is a forest, i.e.* $G(\mathcal{M})$ *contains no cycles.*

Antichain Construction



When does that apply?

Question

When is a class $\mathcal{C}$ (a subset of) a monotone grid class?

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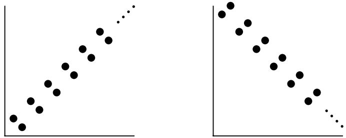
Answer [Vatter]

A class $\mathcal{C}$ is monotone griddable if and only if it contains neither the classes $\oplus 21$ nor $\ominus 12$.

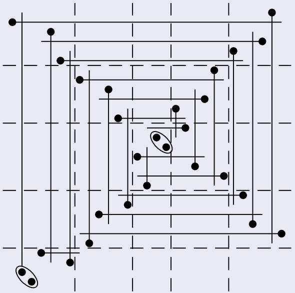


Two is too many

Theorem

A grid class whose graph has a component containing two or more non-monotone-griddable classes is not pwo.

Proof.



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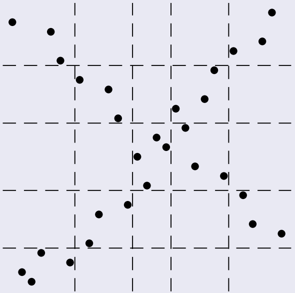

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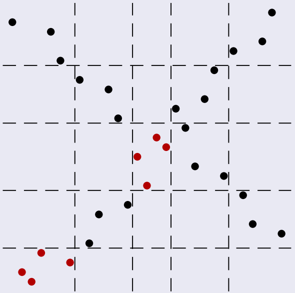

- Antichain element.
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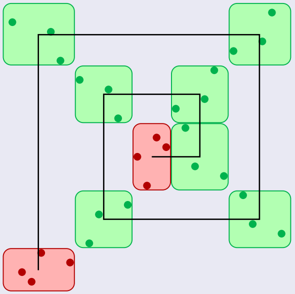

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- Graph is a path.

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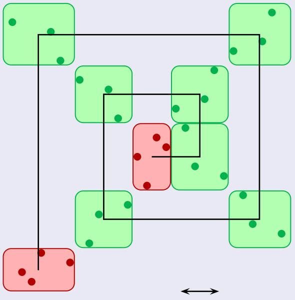

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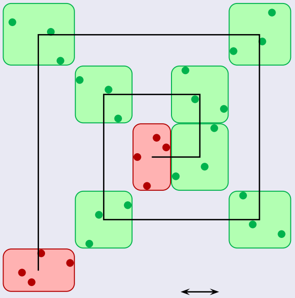

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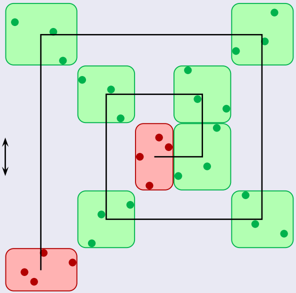

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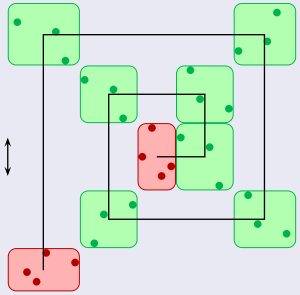

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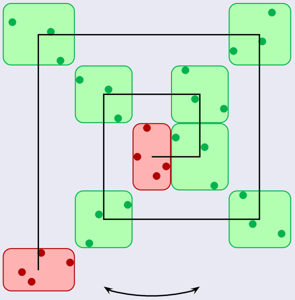

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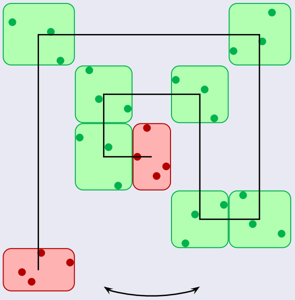

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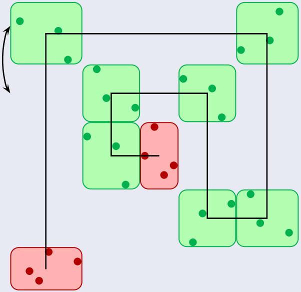

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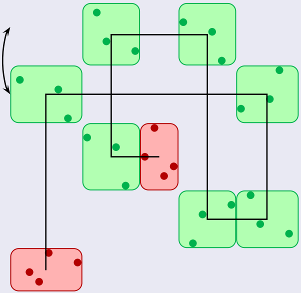

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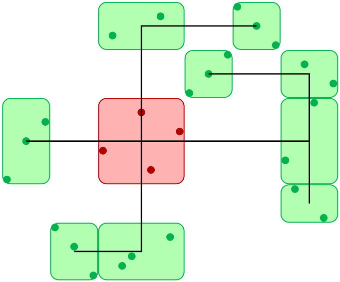
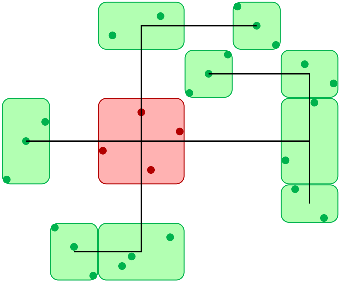
Just one non-monotone

- Bad cell contains only finitely many "simple permutations". ▸ Huh?



Just one non-monotone

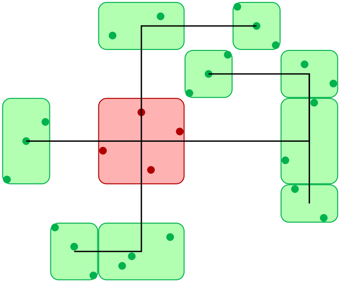
- Bad cell contains only finitely many "simple permutations".
- Form a rooted tree on the red cell, and use Higman's Theorem.



Just one non-monotone

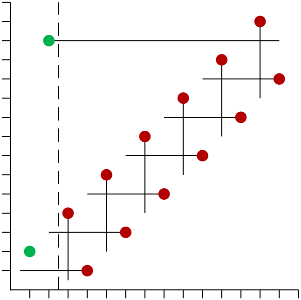
Theorem

Let $\mathcal{M}$ be a gridding matrix for which each component is a forest and contains at most one non-monotone cell. If every non-monotone cell contains only finitely many simple permutations, then $\mathrm{Grid}(\mathcal{M})$ is pwo.



But sometimes one is too much...

- One cell containing arbitrarily long increasing oscillations next to a monotone cell is bad...



Summary

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- To-do: one non-monotone but infinitely many simples (some antichains known).

Question

Can we decide whether a permutation class given by a finite basis is pwo?

- We're closer to answering this, but still some way off.

Thanks!

Appendix: Simple Permutations

- A simple permutation: bounding box around some points is always separated by some other point.

Back

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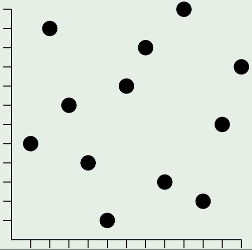
Example



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Appendix: Simple Permutations

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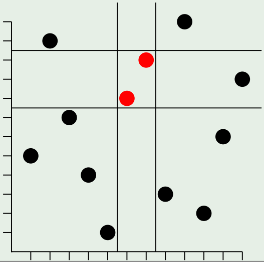
Example



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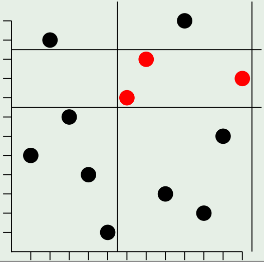
Example



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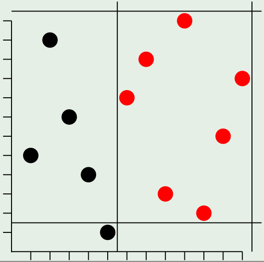
Example



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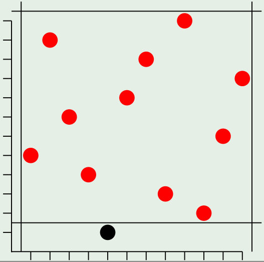
Example



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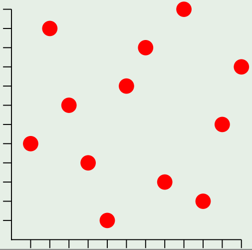
Example



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